



# 武汉大学 人工智能与软件工程暑期学校

2020-08



演讲人：何 强

Qiang HE received his first Ph. D. degree from Swinburne University of Technology (SUT), Australia, in 2009 and his second Ph. D. degree from Huazhong University of Science and Technology (HUST), China, in 2010. He is currently an Associate Professor at Swinburne University of Technology. His major research interests include edge computing, software engineering, service-oriented computing and cloud computing. He has published 110+ papers at top venues, e.g., TPDS, TSE, TSC, TCC, TBD, JPDC, ICSE, WWW, ICDE, IJCAI, ICWS, ICSOC and CLOUD. He is recipient of the Best Student Paper Awards at SCC2018, ICWS2017, ICSOC2019, and the Best Paper Awards at ENASE2020 ICSOC2018.



Best Research Paper ICSOC2018, CORE A

# Optimal Edge User Allocation in Edge Computing with Variable Sized Vector Bin Packing

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- *About Us*
- Background
- Problem Identification
- Problem Statement
- Our Solution
- Experiments



# Swinburne University of Technology

45

top young universities  
in the world by QS  
World University  
Rankings 2021



62

2020 Times Higher  
Education Young  
University Rankings



327

QS World University  
Rankings 2021



129

Computer Science &  
Engineering ranked 129  
over the world by  
Shanghai Ranking



# Our Team

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## Introduction

Our team has 6 (+3 ) PhD students



## Main Research Topics

- ✓ Edge Resources (Since 2018)
  - ✓ Edge User Allocation
  - ✓ Edge Data Caching
  - ✓ Edge Server Placement
  - ✓ Edge Demand Response
- ✓ Edge Security (Since 2019)
  - ✓ Edge Data Integrity
  - ✓ Edge Blockchain
- ✓ Edge AI (Since 2020)



## Major Awards

- ✓ Best Research Paper ENASE2020 (CORE B)
- ✓ Best Research Paper ICSOC2018 (CCF B, CORE A)
- ✓ Best Student Paper ICWS2019 (CCF B , CORE A)
- ✓ Best Student Paper ICSOC2018 (CCF B , CORE A)
- ✓ Best Student Paper ICWS2017 (CCF B , CORE A)
- ✓ ...



# Our Publications on Edge Computing (Since 2018)

Topic	Paper Title	Venue	Year	Rank	
Edge User Allocation	Optimal Edge User Allocation in Edge Computing with Variable Sized Vector Bin Packing	ICSOC	2018	CCF B	CORE A
	Edge User Allocation with Dynamic Quality of Service	ICSOC	2019	CCF B	CORE A
	A Game-theoretical Approach for User Allocation in Edge Computing Environment	TPDS	2020	CCF A	CORE A*
	Quality of Experience-Aware User Allocation in Edge Computing Systems: A Potential Game	ICDCS	2020	CCF B	CORE A
	Cost-Effective App User Allocation in an Edge Computing Environment	TCC	2020	JCR Q1	
	QoE-aware User Allocation in Edge Computing Systems with Dynamic QoS	FGCS	2020	JCR Q1	
	Interference-aware SaaS User Allocation Game for Edge Computing	TCC	2020	JCR Q1	
Edge Data Caching	Online Collaborative Data Caching in Edge Computing	TPDS	2020	CCF A	CORE A*
	Cost-Effective App Data Distribution in Edge Computing	TPDS	2020	CCF A	CORE A*
	Graph-based Optimal Data Caching in Edge Computing	ICSOC	2019	CCF B	CORE A
	Budgeted Data Caching based on k-Median in Mobile Edge Computing	ICWS	2020	CCF B	CORE A
	Graph-based Data Caching Optimization in Edge Computing	FGCS	2020	JCR Q1	



# My Footprints



# My Publications and Awards

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- **Journal Papers (54)**

- **28 x ACM/IEEE Transactions**
- **2 x TPDS, 5 x TSE, 8 x TSC, 5 x TCC, 5 x TBD**

- **Conference Papers (58)**

- **CCF A:** IJCAI, ICDE, WWW, ICSE
- **CCF B:** 8 x ICSOC, 12 x ICWS, 8 x SCC, ICDCS, ICDM, AAMAS

## Major Awards

- **Best Paper Award**, ENASE2020
- **Best Student Paper Award**, ICWS2019 (CCF B)
- **Best Paper Award**, ICSOC2018 (CCF B).
- **Best Student Paper Award**, SCC2018
- **Best Student Paper Award**, ICWS2017 (CCF B)
- **FSET MCR Award**, 2020
- **FSET MCR Award**, 2018



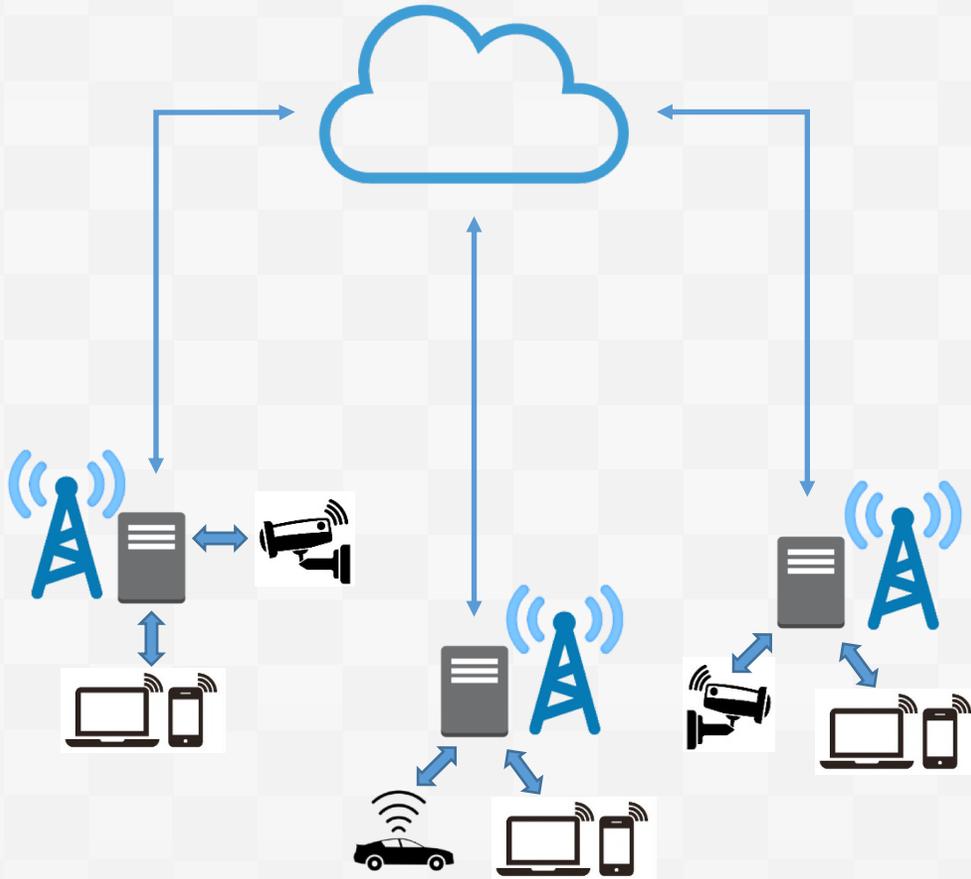
# Outlines

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# What is edge computing and why?



## Advantages

- Computation offloading
- Energy saving
- Low latency

## Applications

- Driverless cars
- Mobile gaming
- Augmented reality
- Remote healthcare
- Smart manufacturing
- ...

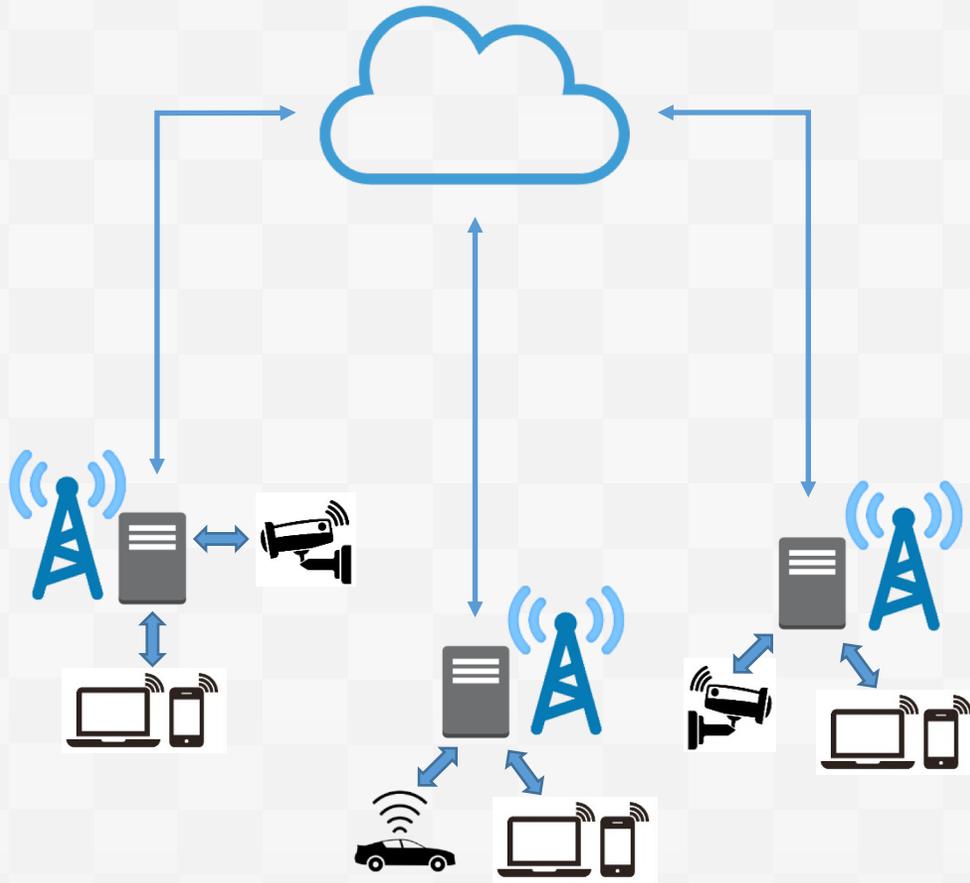
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# Current Research Problem and Perspectives



## Computation Offloading

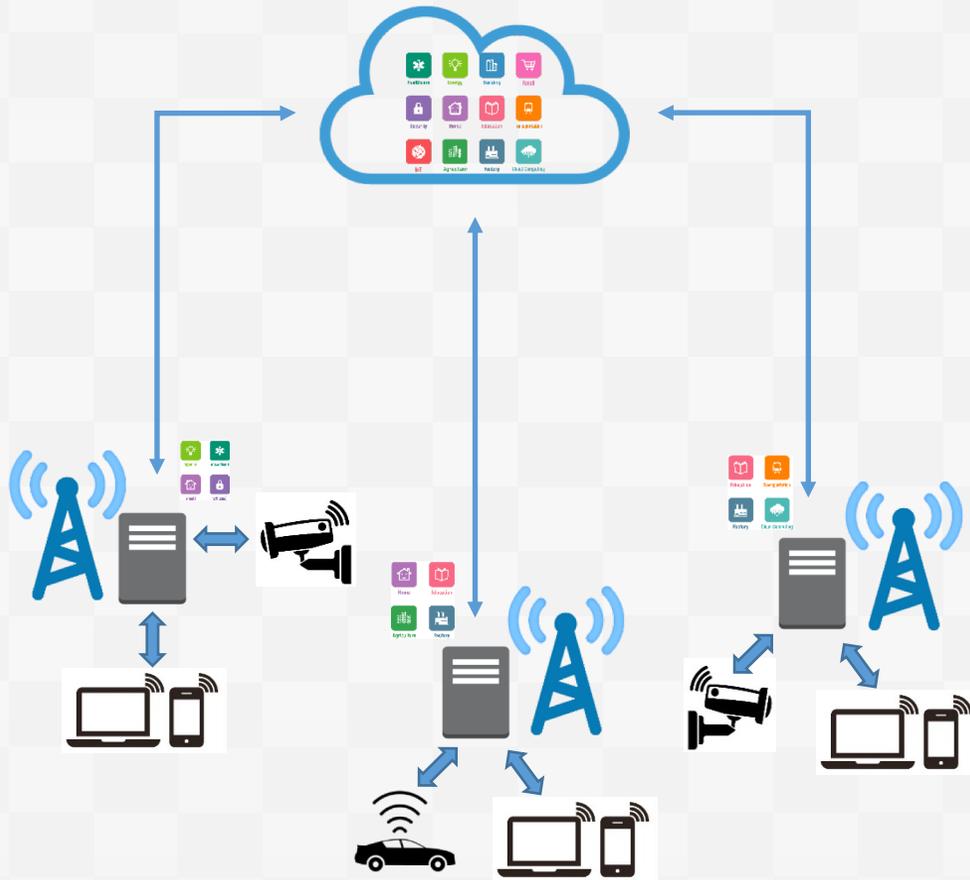
### Mobile/IoT users and devices

- Energy efficiency
- Latency

### Edge infrastructure provider

- Network throughput
- Workload balance

# What about app vendors?



## 💡 App vendor (service provider, content provider)

- Important stakeholder
- Major users of edge infrastructure

## 💡 Major concerns

- Benefit by serving app users
- Cost by hiring computing resources on edge servers

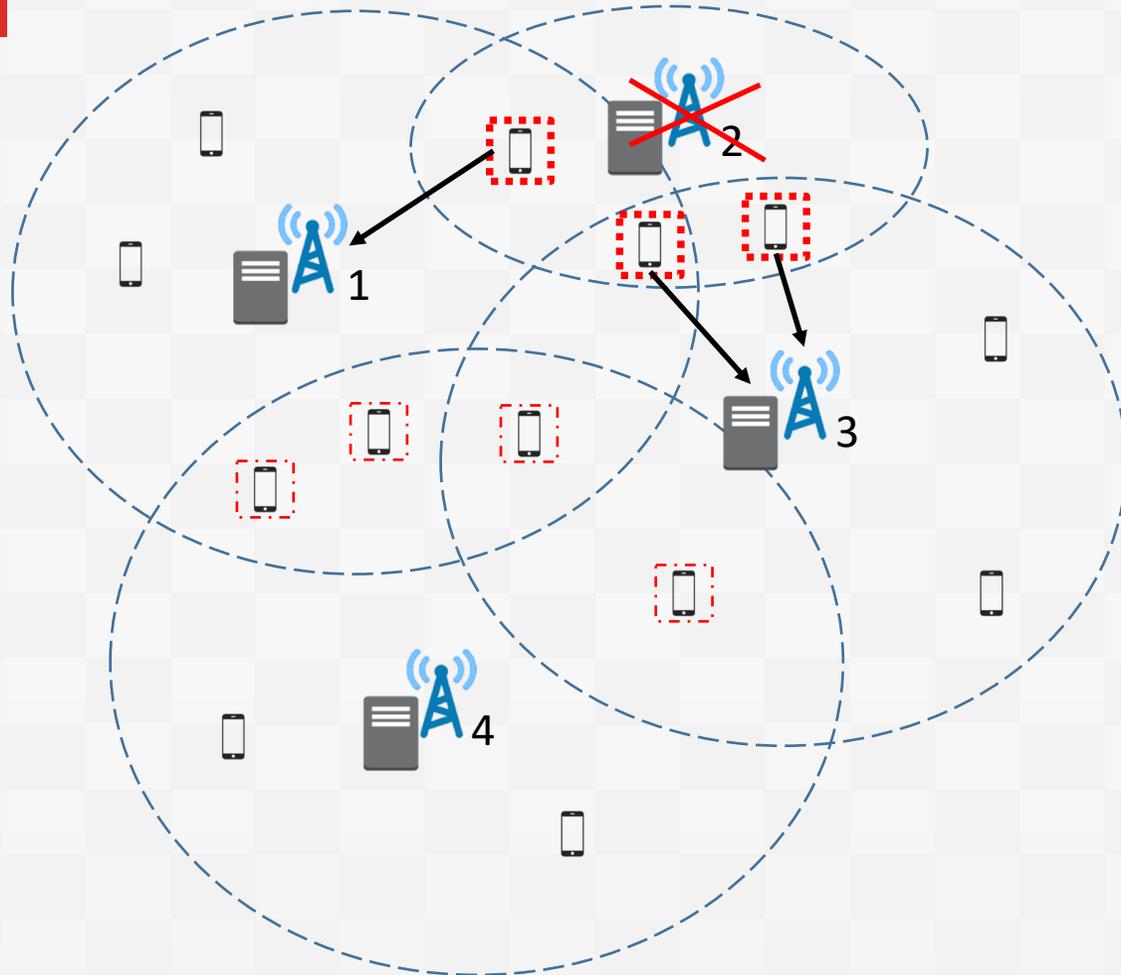
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# Edge User Allocation Problem



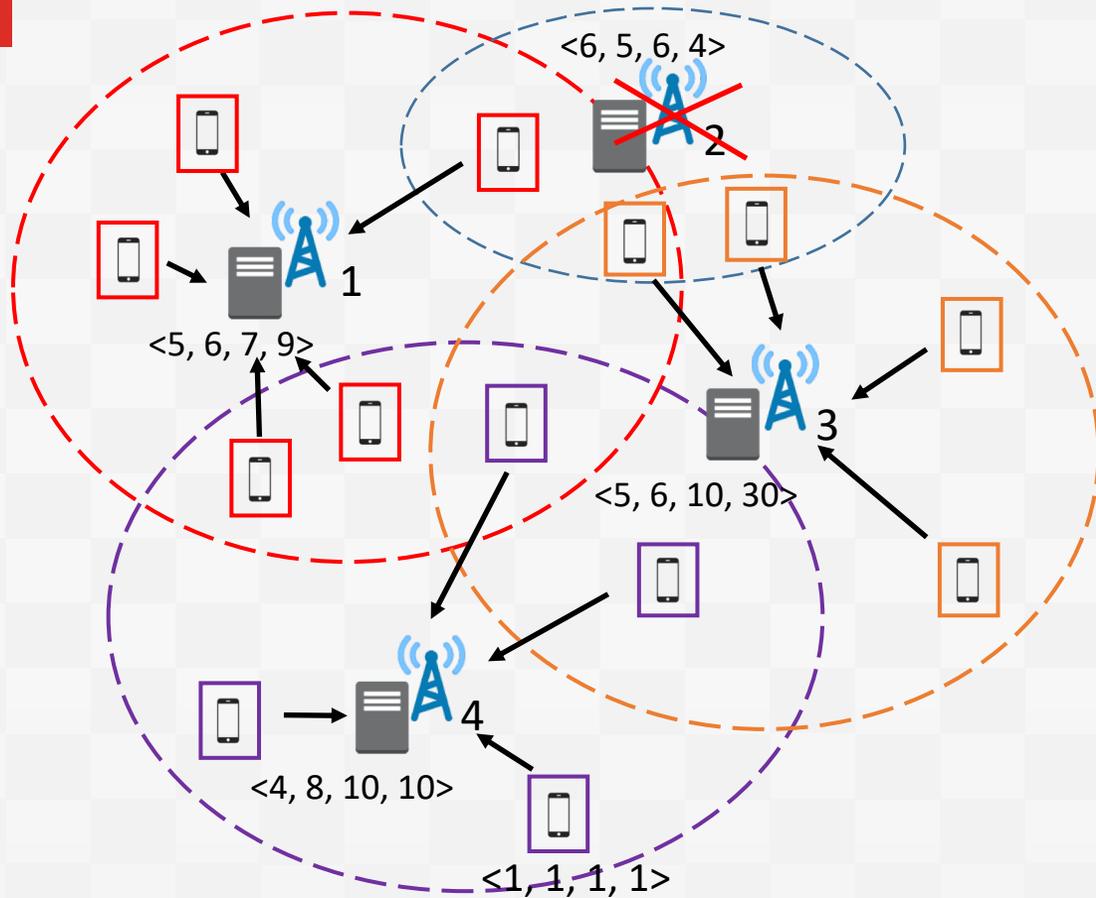
## Objectives

- Maximize number of app users allocated
- Minimize number of edge servers used

## Proximity constraint

- An edge server can only serve users within its coverage

# Edge User Allocation Problem



Capacity vector: <CPU, RAM, VRAM, Bandwidth>

## Objectives

- Maximize number of app users allocated
- Minimize number of edge servers used

## Proximity constraint

- An edge server can only serve users within its coverage

## Capacity constraint

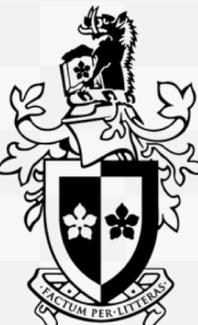
- Demands of users allocated to an edge server must not exceed its remaining capacities



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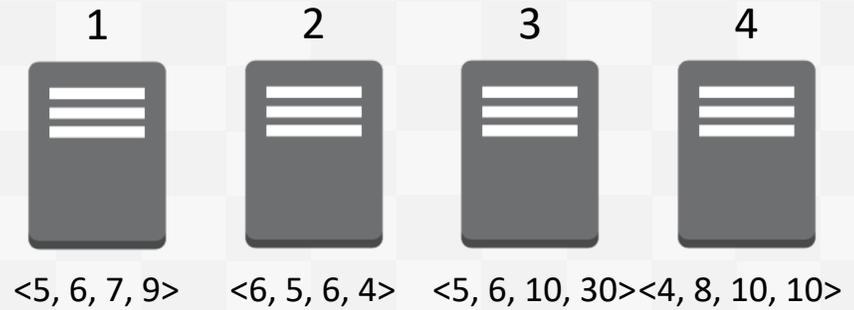
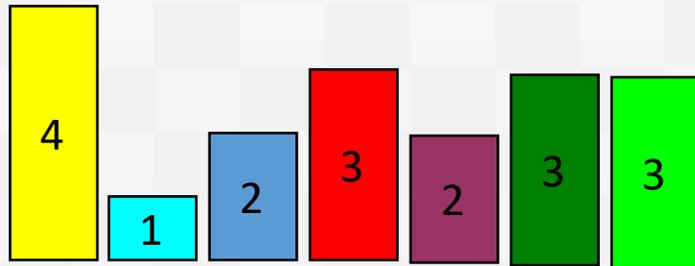
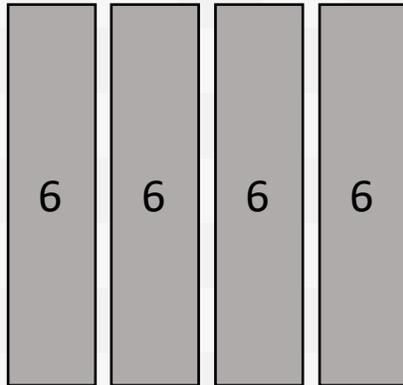


# Problem Modelling

Bin Packing Problem (*NP*-hard)



**Variable Sized Vector** Bin Packing Problem (*NP*-hard)



<1, 1, 1, 1>

Capacity vector: <CPU, RAM, VRAM, Bandwidth>



# Integer Programming Optimization

## ● Optimization objectives:

- Maximize the number of allocated users
- Minimize the number of used servers

## ● Constraints:

- Proximity constraint
- Capacity constraint

## ● Solver: IBM CPLEX Optimizer



Notation	Description
$S = \{s_1, s_2, \dots, s_i\}$	finite set of edge server $s_i$ , where $i = 1, 2, \dots, m$
$C_i = \langle C_i^1, C_i^2, \dots, C_i^d \rangle$	$d$ -dimensional vector with each dimension $C_i^k$ being a resource type, such as CPU utilization or disk I/O, representing the remaining capacity of an edge server $s_i$ , $k \in \{1, 2, \dots, d\}$
$U = \{u_1, u_2, \dots, u_j\}$	finite set of user $u_j$ , where $j = 1, 2, \dots, n$
$w_j = \langle w_j^1, w_j^2, \dots, w_j^d \rangle$	$d$ -dimensional vector representing the size of the workload incurred by user $u_j$ . Each vector component $w_j^k$ is a resource type, $k \in \{1, 2, \dots, d\}$
$U(s_i)$	set of users allocated to server $s_i$ . $U(s_i) \subset U$
$d_{ij}$	geographical distance between server $s_i$ and user $u_j$
$cov(s_i)$	coverage radius of server $s_i$

$$\text{maximize } \sum_{j=1}^n \sum_{i=1}^m x_{ij} \quad (1)$$

$$\text{minimize } E = \sum_{i=1}^m y_i \quad (2)$$

subject to:

$$\sum_{j=1}^n w_j^k x_{ij} \leq C_i^k y_i, \forall i \in \{1, \dots, m\}; \forall k \in \{1, \dots, d\} \quad (3)$$

$$d_{ij} \leq cov(s_i), \forall i \in \{1, \dots, m\}; \forall j \in \{1, \dots, n\} \quad (4)$$

$$\sum_{i=1}^m x_{ij} \leq 1, \forall j \in \{1, \dots, n\} \quad (5)$$

$$y_i \in \{0, 1\}, \forall i \in \{1, \dots, m\} \quad (6)$$

$$x_{ij} \in \{0, 1\}, \forall i \in \{1, \dots, m\}; \forall j \in \{1, \dots, n\} \quad (7)$$

where:

$y_i = 1$  if server  $s_i$  is hired.

$x_{ij} = 1$  if user  $u_j$  is allocated to server  $s_i$ .

$cov(s_i)$  is provided by edge computing providers.



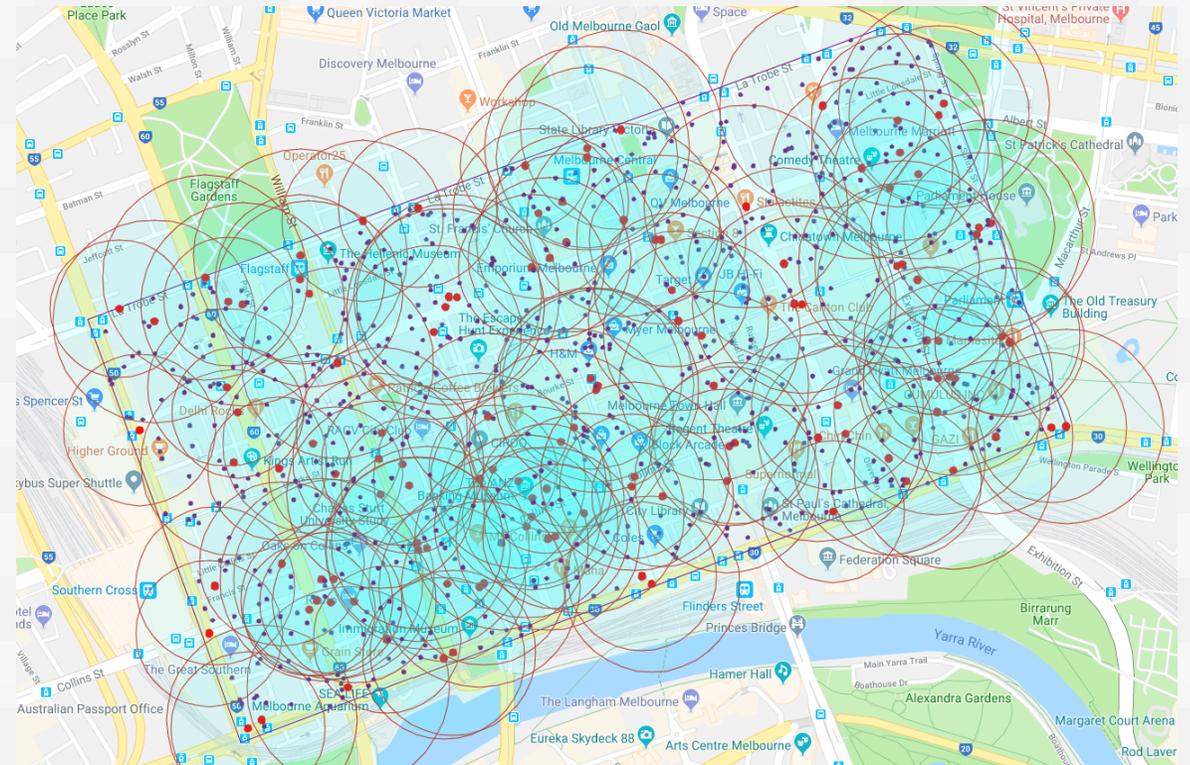
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- **Experiments**



- **Area:** Simulated Melbourne CBD area
- **Edge servers:**
  - 125 Telstra base stations in the CBD area.
  - Coverage: 450-750m.
- **End-users:**
  - 550 users in the CBD area.



Datasets: <https://github.com/swinedge/eua-dataset> or <https://sites.google.com/site/heqiang/eua-repository>, containing **95,562** base stations in Australia and **~131,000** users.



# Experiment Settings

## ● Comparing approaches

- **Random**: Randomly allocates end-users
- **Greedy**: Always allocates the most end-users to an edge server.

## ● Parameters settings

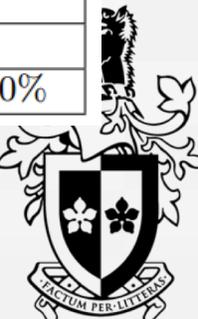
- **Number of end-users**: randomly select 4, 8, 16, ..., 512 users.
- **Number of edge servers**: 10%, 20%, ..., 100% of the number of the servers are available.
- **Remaining server capacity**: 100%, 150%, ..., 300% of the combined user workload are available.

## ● Performance metrics

- **Percentage of users allocated**
- **Percentage of used edge servers**
- **Execution time (CPU time)**

Table 2: Experiment Settings

Factor	Number of users	Number of servers	Remaining server capacity
Set #1	4, 8, ..., 512	100%	300%
Set #2	512	10%, 20%, ..., 100%	300%
Set #3	512	100%	100%, 150%, ..., 300%



# Experiment Results

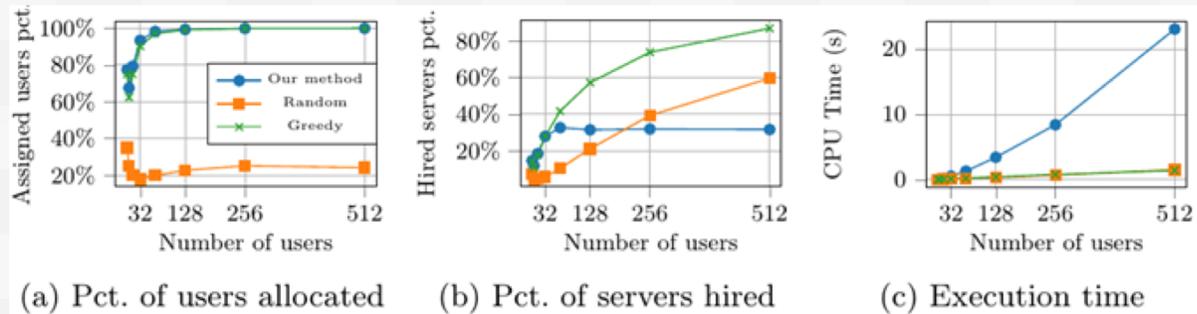


Fig. 1 Results of experiment set #1

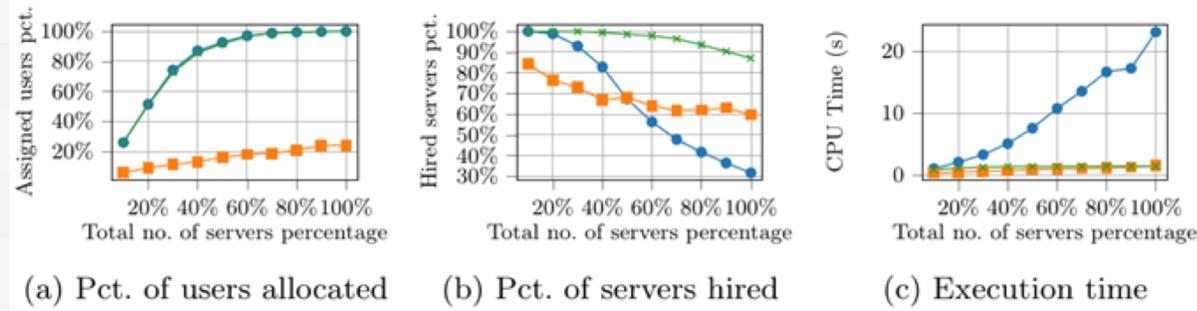


Fig. 2 Results of experiment set #2

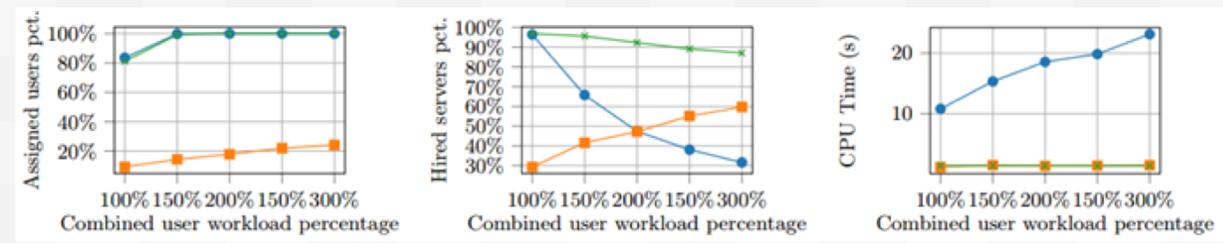


Fig. 3 Results of experiment set #3



# Our Major Contributions

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- ✓ Identify the new edge user allocation (EUA) problem
- ✓ Model EUA as Variable Sized Vector Bin Packing Problem
- ✓ Propose a solution based on Integer Programming



# More Work on Edge User Allocation

- Edge User Allocation with **Dynamic Quality of Service** (ICSOC2019, CCF B)
- A **Game-theoretical** Approach for User Allocation in Edge Computing Environment (TPDS2019, CCF A, JCR Q1)
- **Quality of Experience-Aware** User Allocation in Edge Computing Systems: A Potential Game (ICDCS2020, CCF B)
- **Cost-Effective** App User Allocation in an Edge Computing Environment (TCC2020, JCR Q1)
- **QoE-aware** User Allocation in Edge Computing Systems with **Dynamic QoS** (FGCS2020, JCR Q1)
- **Interference-aware** SaaS User Allocation Game for Edge Computing (TCC2020, JCR Q1)





***THANK  
YOU***



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